

# Algorithms and Data Structures

Summer Semester 2025

For discussion on Wednesday, May 21, 2025

1. (GTG Exercise R-5.3) Modify the experiment from

[https://github.com/mjwestcott/Goodrich/blob/master/ch05/experiment\\_list\\_size.py](https://github.com/mjwestcott/Goodrich/blob/master/ch05/experiment_list_size.py)

in order to demonstrate that Python's list class occasionally shrinks the size of its underlying array when elements are popped from a list.

2. (GTG Exercise R-5.7) Let  $A$  be an array of size  $n \geq 2$  containing integers from 1 to  $n - 1$ , inclusive, with exactly one repeated. Describe a fast algorithm for finding the integer in  $A$  that is repeated.
3. (GTG Exercise C-5.15) Consider an implementation of a dynamic array, but instead of copying the elements into an array of double the size (that is, from  $N$  to  $2N$ ) when its capacity is reached, we copy the elements into an array with  $\lceil N/4 \rceil$  additional cells, going from capacity  $N$  to capacity  $N + \lceil N/4 \rceil$ . Prove that performing a sequence of  $n$  append operations still runs in  $O(n)$  time in this case.
4. (From GTG Exercise C-5.19) Consider a shrink strategy in which an array of capacity  $N$  is resized to capacity precisely that of twice the number of elements any time the number of elements in the array goes strictly below  $N/4$ . Give a formal proof that any sequence of  $n$  append or pop operations on an initially empty dynamic array takes  $O(n)$  time.
5. (GTG Exercise C-5.20) Consider a variant shrink strategy in which a dynamic array of capacity  $N$  is resized to capacity precisely that of the number of elements, any time the number of elements in the array goes strictly below  $N/2$ . Show that there exists a sequence of  $n$  operations that requires  $\Omega(n^2)$  time to execute.