

1. A connected graph  $G$  is called 2-connected if it remains connected after removal of any one of its vertices.

(a) Give an example of a graph that is connected, but not 2-connected. (5)

(b) Show that a graph with at least three vertices is 2-connected if and only if every pair of vertices lies in a cycle. (5+5)

a) Any tree with more than two vertices.

b) " $\Rightarrow$ ": Fix any two distinct vertices  $a$  and  $b$ .

Since  $G$  is 2-connected,  $a$  must be part of some cycle  $C$ .

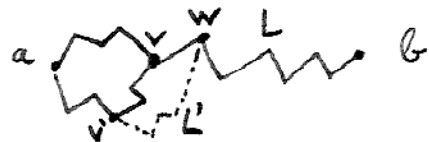
If  $b \in C$ , we are done. Otherwise, we extend  $C$  along a simple chain  $L$  from a vertex  $v \in C$  to  $b$  (which must exist as  $G$  is connected) as follows: Let  $w \in L$  be the

vertex neighboring  $C$ . As  $G$  is two-connected, there is a simple chain  $L'$  from  $w$  to some  $v' \in C$  avoiding  $v$ .

Then the concatenation of  $C$  without the piece from  $v$  to  $v'$ ,

the edge  $(v, w)$ , and  $L'$  are again a cycle which now contains  $w$ . Continuing, we reach  $w = b$  in a finite

number of steps along  $L$ .



" $\Leftarrow$ ": Fix  $v \in G$  and let  $a, b \in G \setminus \{v\}$  be arbitrary. Since there is a cycle containing  $a$  and  $b$ , there is at least one path between  $a$  and  $b$  avoiding  $v$ .

2. If  $A$  and  $B$  are points in the plane, let  $U_{AB}$  denote the glide reflection along the line  $AB$  which maps  $A$  to  $B$ .

(a) Given a rectangle with vertices  $A, B, C,$  and  $D$ , show that

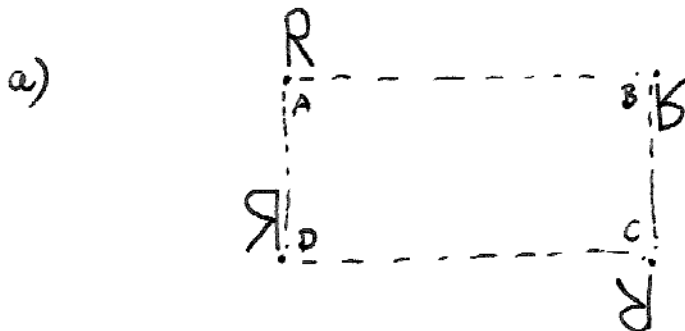
$$U_{CD} \circ U_{BC} \circ U_{AB} \circ U_{DA} = e.$$

(In other words, "gliding around the a rectangle" is the identity.)

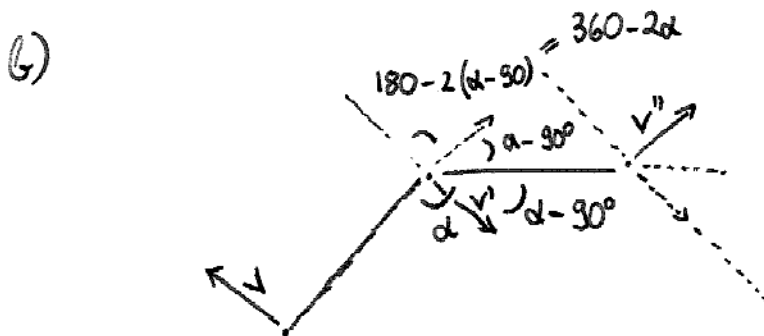
(b) Find a condition for a more general quadrilateral that ensures that gliding around it also results in the identity.

*Hint:* If the angle between  $AB$  and  $BC$  is  $\alpha$ , by which angle does  $U_{BC} \circ U_{AB}$  rotate a vector?

(5+5)



Following the orientation of the letter  $R$  as it is mapped around the rectangle proves the claim.



$$v' = U_{AB} v$$

$$v'' = U_{BC} v'$$

So the first two glide reflections rotate a vector by  $2\alpha$  (ignoring the sign), so if  $\beta$  is the angle opposite  $\alpha$ , we need  $2\alpha + 2\beta = 360^\circ$  or  $\alpha + \beta = 180^\circ$ .

3. (a) Characterize the group of motions of the line, i.e., the group of maps  $\phi: \mathbb{R} \rightarrow \mathbb{R}$  which preserve the distance between points.
- (b) Prove that the set of matrices

$$G = \left\{ \begin{pmatrix} \pm 1 & \lambda \\ 0 & 1 \end{pmatrix} : \lambda \in \mathbb{R} \right\}$$

is a group with respect to the usual matrix multiplication. Is it Abelian?

- (c) Show that the group of motions of the line is isomorphic to  $G$ .

*Hint:* Show that the set

$$L = \left\{ \begin{pmatrix} x \\ 1 \end{pmatrix} : x \in \mathbb{R} \right\}$$

is invariant under  $G$ .

(5+5+5)

- a) If  $y = \phi(x)$ , then either  $\phi(x+\delta) = y+\delta$  or  $\phi(x+\delta) = y-\delta$  as  $\phi$  must preserve distances. In the first case,  $\phi$  is a translation, in the second case it is a composition of translation and reflection.

Thus, the group is generated by all translations and one reflection.

- b) By direct computation, we see that

$$\begin{pmatrix} 1 & \lambda \\ 0 & 1 \end{pmatrix}^{-1} = \begin{pmatrix} 1 & -\lambda \\ 0 & 1 \end{pmatrix}$$

and

$$\begin{pmatrix} -1 & \lambda \\ 0 & 1 \end{pmatrix}^{-1} = \begin{pmatrix} -1 & \lambda \\ 0 & 1 \end{pmatrix}$$

Moreover,

$$\begin{pmatrix} 1 & \lambda \\ 0 & 1 \end{pmatrix} \begin{pmatrix} \pm 1 & \mu \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} \pm 1 & \mu + \lambda \\ 0 & 1 \end{pmatrix}$$

and

$$\begin{pmatrix} -1 & \lambda \\ 0 & 1 \end{pmatrix} \begin{pmatrix} \pm 1 & \mu \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} \mp 1 & \lambda - \mu \\ 0 & 1 \end{pmatrix}$$

I.e.,  $G$  is closed under inversion and matrix multiplication. Thus, it is a subgroup of the group of invertible  $2 \times 2$  matrices, hence a group.

c)

$$\begin{pmatrix} \pm 1 & \lambda \\ 0 & 1 \end{pmatrix} \begin{pmatrix} x \\ 1 \end{pmatrix} = \begin{pmatrix} \pm x + \lambda \\ 1 \end{pmatrix}$$

Thus, if we identify  $x \in \mathbb{R}$  with  $\begin{pmatrix} x \\ 1 \end{pmatrix} \in \mathbb{R}^2$ , then the matrix  $\begin{pmatrix} \pm 1 & \lambda \\ 0 & 1 \end{pmatrix}$  corresponds to the motion  $\phi(x) = \pm x + \lambda$ ; by (a) any motion on the line is of this form.

4. A tailor makes jackets and pants. There is enough demand that she sells everything she produces. It takes an hour to make a jacket and half an hour to make a pair of pants. She can spare 10 hours per week for sewing and has a long-running supply contract that provides cloth for 15 pieces per week altogether. The profit on a pair of pants is EUR 15 and the profit on a jacket is EUR 20. How many pieces of each per week should she produce to maximize profit? (10)

Let  $x$  denote the number of jackets per week,  
 $y$  denote the number of pants per week.

Then  $x + y \leq 15$   
 $x + \frac{1}{2}y \leq 10$

and she ought to maximize  $z = 20x + 15y$

In standard form:

$$x + y + s_1 = 15$$

$$x + \frac{1}{2}y + s_2 = 10$$

$$x, y, s_1, s_2 \geq 0$$

minimize  $J = -20x - 15y$ .

Let's use the simplex method (graphical is easy as well):

		↓ $s_1$	↓ $s_2$	
$x$	$y$			
1	1	1	0	15
1	$\frac{1}{2}$	0	1	10
-20	-15	0	0	0

$x$  is entering

$s_2$  is leaving

→

↓ $x$	$y$	↓ $s_1$	$s_2$	
0	$\frac{1}{2}$	1	-1	5
1	$\frac{1}{2}$	0	1	10
0	-5	0	20	200

$y$  is entering

$s_1$  is leaving

→

$x$	$y$	$s_1$	$s_2$	
0	1	2	-2	10
1	0	-1	3	5
0	0	10	10	250

the method terminates here.

⇒ She should make  $x=5$  jackets and  $y=10$  pants  
at a maximum profit of EUR 250.

5. Suppose each of the following tableaus occurs in the course of performing the simplex algorithm on a linear programming problem in standard form.

(a) 

$x_1$	$x_2$	$x_3$	$x_4$	
0	-1	1	-1	2
1	0	0	2	3
0	-1	0	3	4

$x_2$  is entering variable, but there is no positive pivot  $\Rightarrow$  unbounded optimal solution

(b) 

$x_1$	$x_2$	$x_3$	$x_4$	
0	0	0	1	-1
0	0	1	0	1
2	1	0	0	10

Solution is  $x_4 = -1$  (no chance to choose different basic variables)  
 $\Rightarrow$  feasible region is empty

(c) 

$x_1$	$x_2$	$x_3$	$x_4$	
2	1	0	1	0
1	0	1	4	3
2	0	0	0	8

Basic variable  $x_2 = 0$   
 $\Rightarrow$  degenerate finite optimal solution.

State, for each case, whether

- The feasible region is empty or nonempty;
- The problem has a finite solution;
- if so, whether the solution is degenerate or nondegenerate.

(10)

6. Let

$$\tilde{v}_k = \frac{1}{N} \sum_{j=0}^{N-1} e^{-ikjh} v_j$$

denote the discrete Fourier transform of the complex numbers  $v_0, \dots, v_{N-1}$ .

Prove the discrete Parseval identity

$$\sum_{k=0}^{N-1} |\tilde{v}_k|^2 = \frac{1}{N} \sum_{j=0}^{N-1} |v_j|^2.$$

(10)

$$\begin{aligned} \sum_{k=0}^{N-1} |\tilde{v}_k|^2 &= \sum_{k=0}^{N-1} \frac{1}{N} \sum_{j=0}^{N-1} e^{-ikjh} v_j \frac{1}{N} \sum_{l=0}^{N-1} e^{iklh} \overline{v_l} \\ &= \frac{1}{N} \sum_{j,l=0}^{N-1} v_j \overline{v_l} \underbrace{\frac{1}{N} \sum_{k=0}^{N-1} e^{ik(l-j)h}}_{= \delta_{lj}^{\text{per}}} \\ &= \frac{1}{N} \sum_{j=0}^{N-1} |v_j|^2 \end{aligned}$$



7. Let  $G$  be a finite Abelian group of order  $N$ , and let  $\chi$  be a character, i.e., a group homomorphism from  $G$  to  $\mathbb{C} \setminus \{0\}$ . Show that  $\chi(a)$  is a root of unity for every  $a \in G$ . (10)

Since  $G$  has order  $N$

$$\chi(a)^N = \chi(Na) = \chi(0) = 1$$

↑  
since  $\chi$  is a group homomorphism  
and with the convention that we  
write  $G$  additively.

$\Rightarrow \chi(a)$  is an  $N$ th root of unity.