

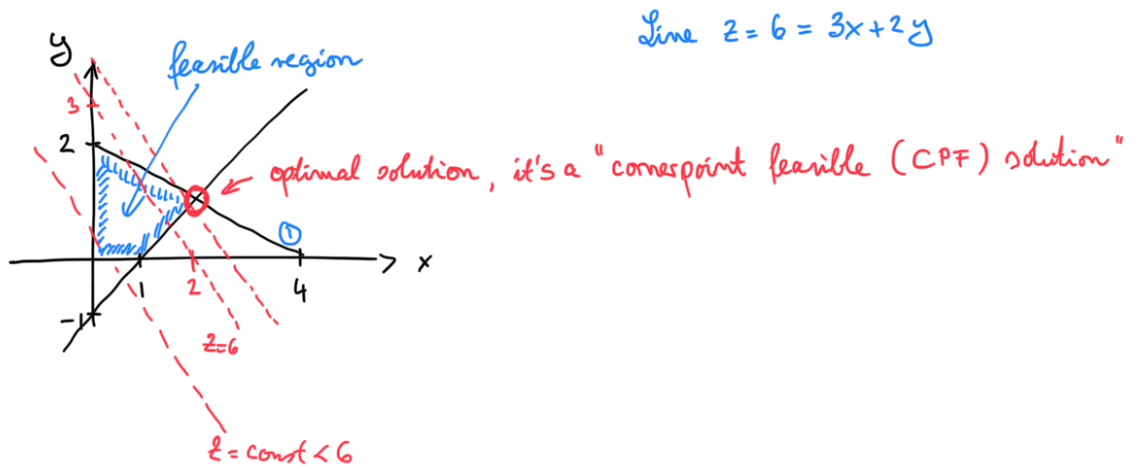
Graphical solution of linear programs

Example 1: maximize $z = 3x + 2y$

subject to $x + 2y \leq 4$ ①

$x - y \leq 1$ ②

$x, y \geq 0$



At the optimal corner point: $x + 2y = 4$
 $x - y = 1$

$$\begin{pmatrix} 1 & 2 & | & 4 \\ 1 & -1 & | & 1 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 2 & | & 4 \\ 0 & -3 & | & -3 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 0 & | & 2 \\ 0 & 1 & | & 1 \end{pmatrix}$$

Solution is $\begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$ with $z = 3 \cdot 2 + 2 \cdot 1 = 8$

Example 2: minimize $z = 6x + 9y$

subject to $3x - y \leq 15$ ①

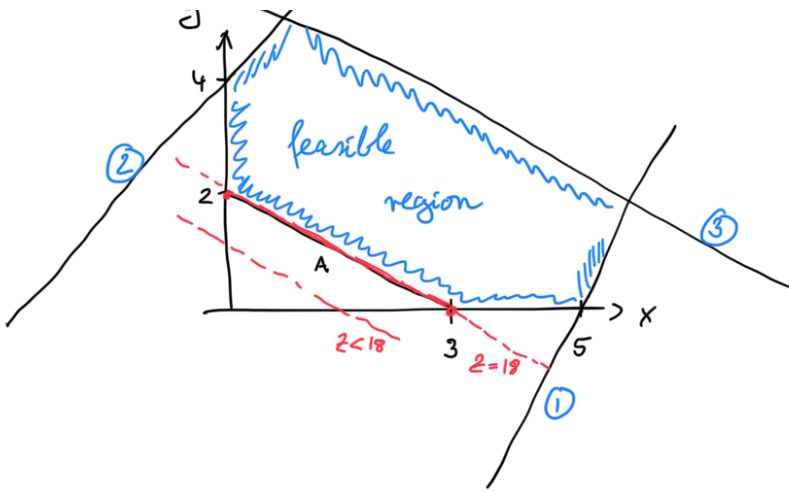
$-x + y \leq 4$ ②

$2x + 5y \leq 27$ ③

$2x + 3y \geq 6$

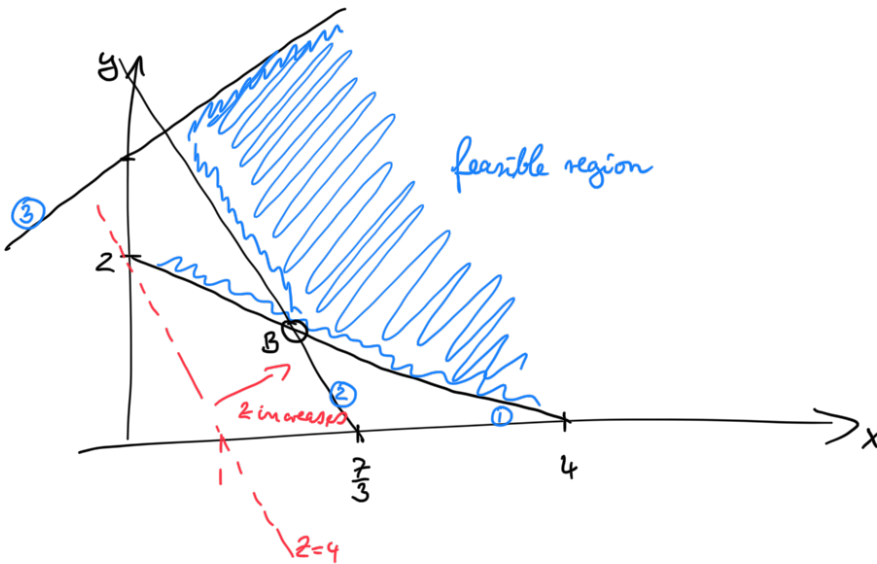
$x, y \geq 0$

$$z = 18 = 6x + 3y$$



Any point on the edge A of the feasible region is an optimal solution
(infinitely many optimal solutions — the problem is degenerate)

Example 3: maximize $z = 4x + 2y$
 subject to $x + 2y \geq 4$ (1)
 $3x + y \geq 7$ (2)
 $-x + 2y \leq 7$ (3)
 $x, y \geq 0$



The feasible region is unbounded and z is increasing in the unbounded direction, so there are feasible solutions, but we cannot find an optimal solution.

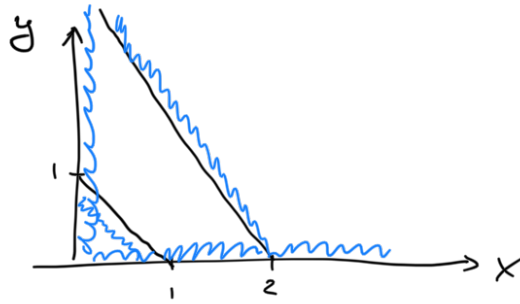
Remark: The minimum of z exists and is taken at corner point B

There: $x + 2y = 4$
 $3x + y = 7$

$$\left(\begin{array}{cc|c} 1 & 2 & 4 \\ 3 & 1 & 7 \end{array} \right) \rightarrow \left(\begin{array}{cc|c} 1 & 2 & 4 \\ 0 & -5 & -5 \end{array} \right) \rightarrow \left(\begin{array}{cc|c} 1 & 0 & 2 \\ 0 & 1 & 1 \end{array} \right)$$

So $\begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$ with $z = 4 \cdot 2 + 2 \cdot 1 = 10$

Example 4: maximize $z = 3x + 4y$
 subject to $x + y \leq 1$
 $2x + y \geq 4$
 $x, y \geq 0$



Here: feasible region is empty
 \Rightarrow no feasible solutions

"the problem is over-constrained"

Prototypical problems:

① Activity analysis problem:

A: set of activities (or products)

R: set of resources (or production facilities)

w_{ij} : workload required from activity $i \in A$ on resource $j \in R$

c_j : available capacity of resource $j \in R$

p_i : profit from performing one unit of activity $i \in A$

Decision variables:

x_i : number of units of activity $i \in A$ to perform

maximize $z = \sum_{i \in A} p_i x_i$
 subject to: $\sum_{i \in A} w_{ij} x_i \leq c_j$ for all $j \in R$
 $x_i \geq 0$ for all $i \in A$

2. The diet problem

Given: F : set of foods

N : set of nutrients

c_i : unit cost of food $i \in F$

r_j : minimum requirements for nutrient $j \in N$

a_{ij} : amount of nutrient $j \in N$ from eating one unit of food $i \in F$

Decision variables:

x_i : # of units of food $i \in F$ to consume

$$\text{minimize } z = \sum_{i \in F} c_i x_i$$

$$\text{subject to } \sum_{i \in F} a_{ij} x_i \geq r_j \quad \text{for every } j \in N$$

$$x_i \geq 0 \quad \text{for } i \in F$$